

# Consensus Protocols II

Murat Osmanoglu

# Consensus Protocols for Blockchain

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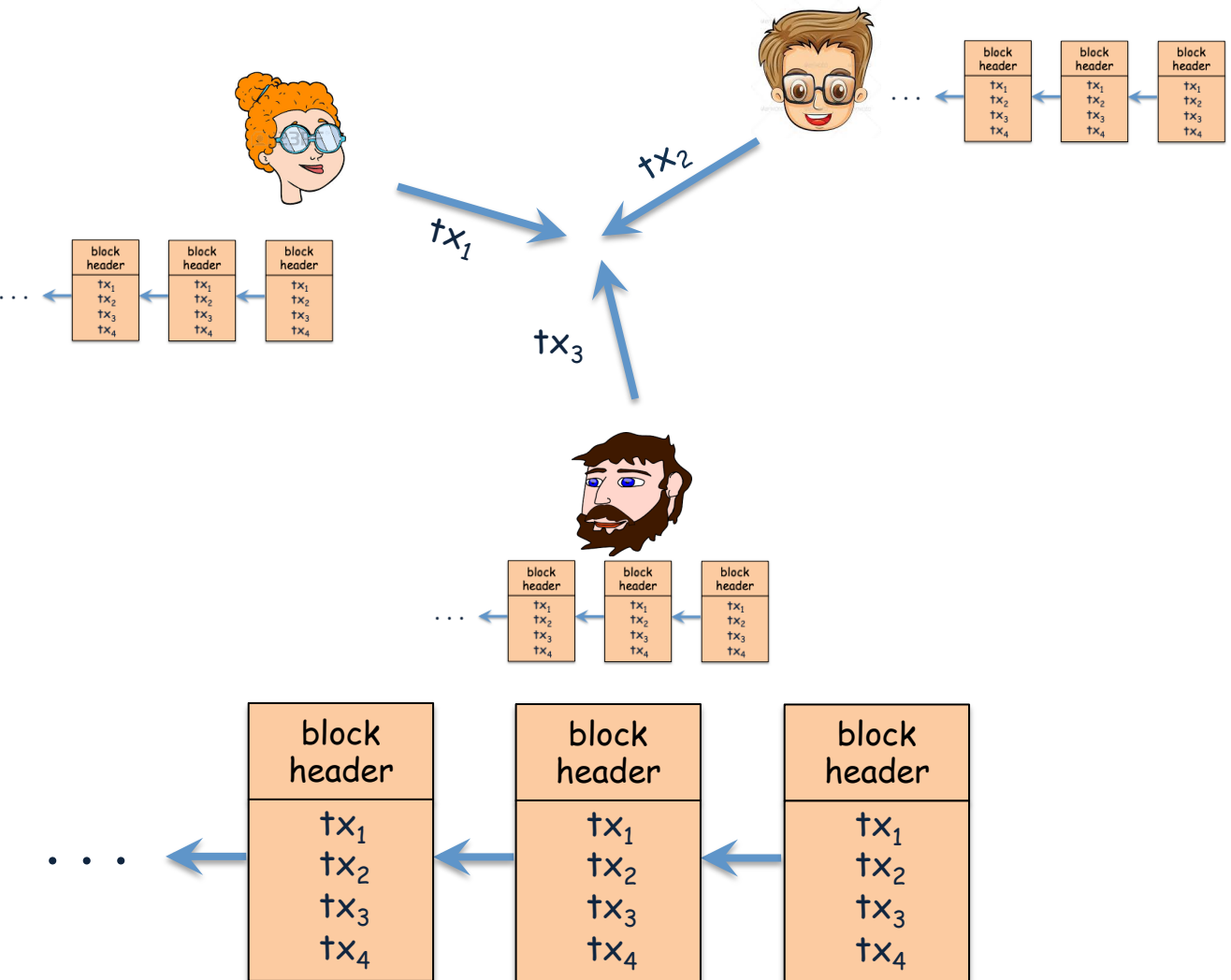
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  - leaders randomly selected with a probability in proportion to some criteria such as its computing power, or its stake
  - newly created blocks appended to the chain without using BFT type voting mechanism
  - block finalization probabilistic

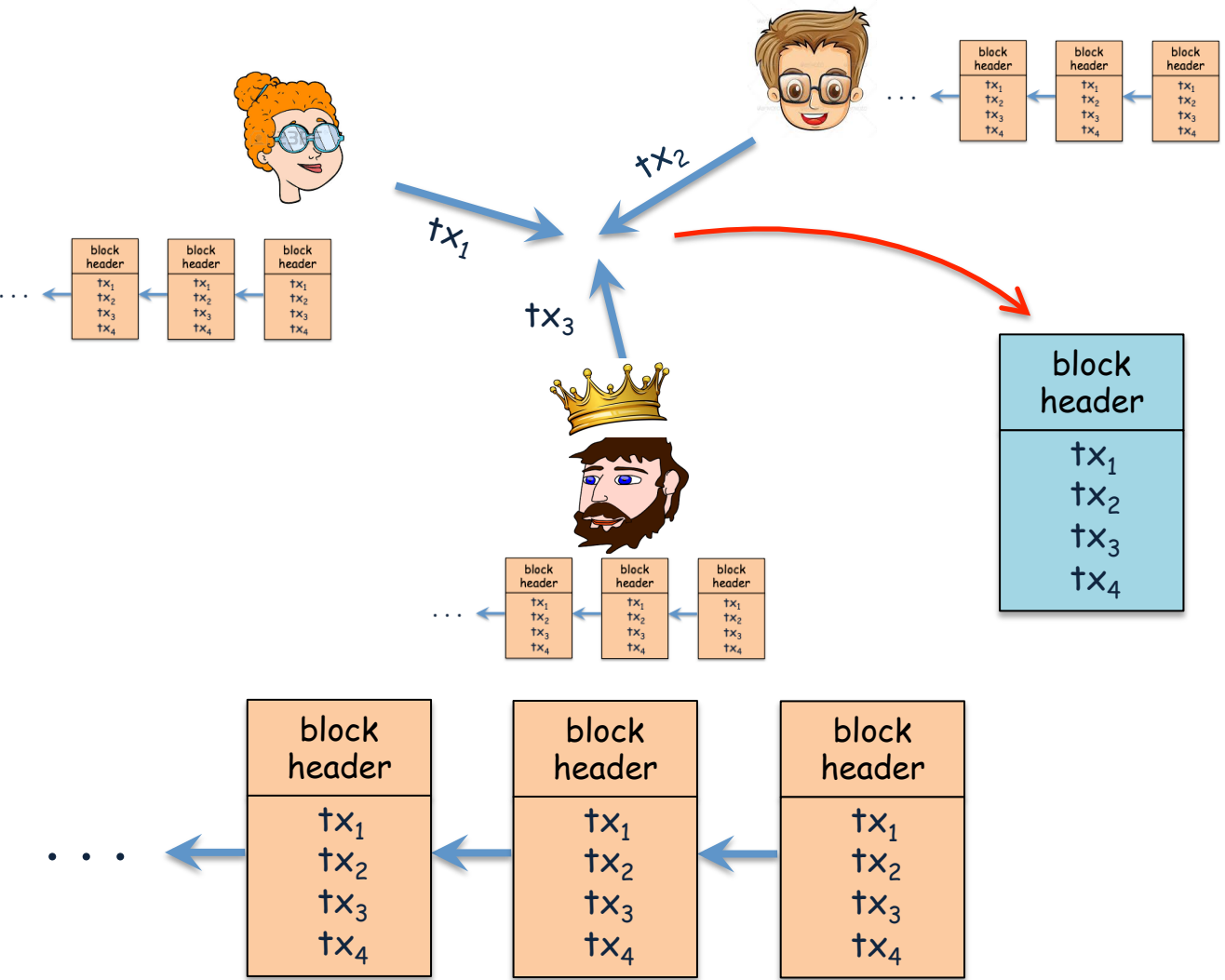
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- voting-based consensus protocol,
  - leaders determined by utilizing simpler methods (round-robin vs.)
  - newly created blocks appended to the chain through BFT type voting mechanism
  - block finalization deterministic

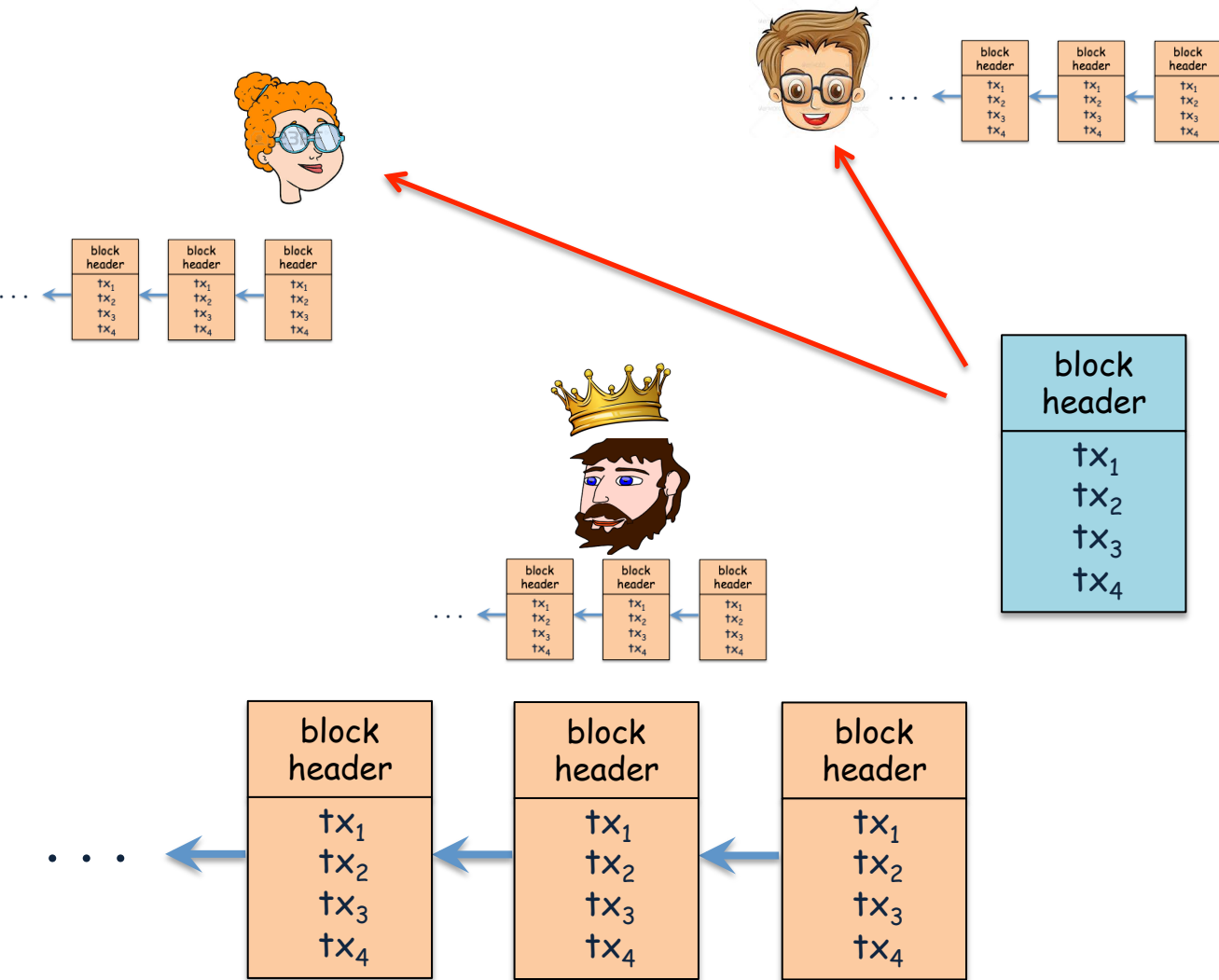
# Blockchain Consensus Protocols-Nakamoto Consensus Protocol (Proof-of-Work Based)



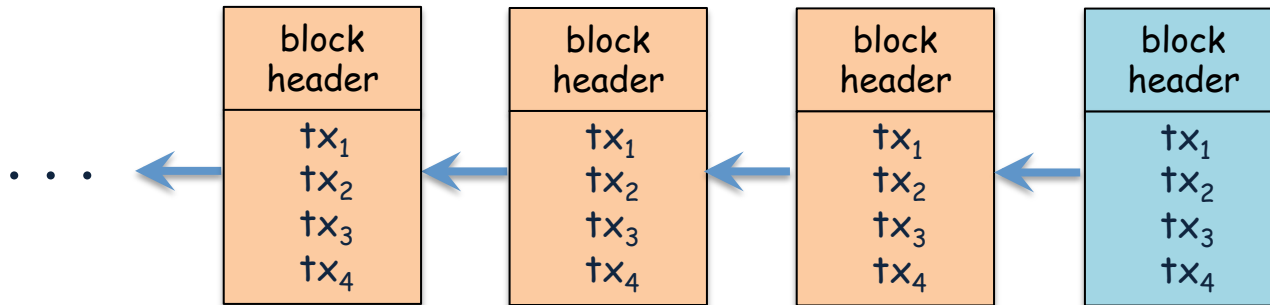
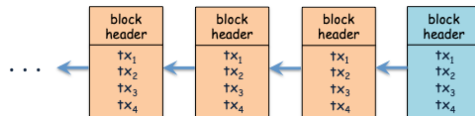
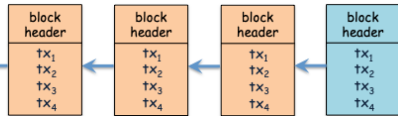
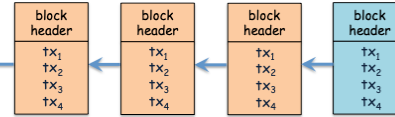
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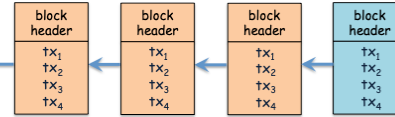
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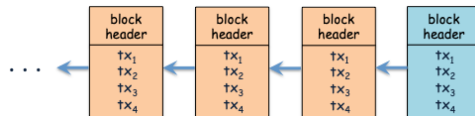
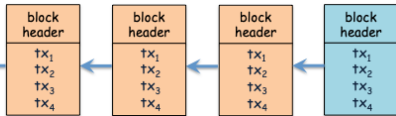
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- validate the signature of miner
- validate the transactions contained in the block



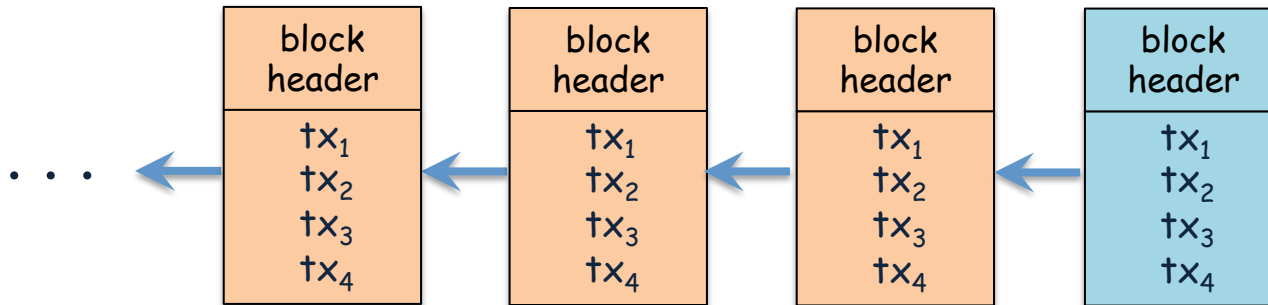
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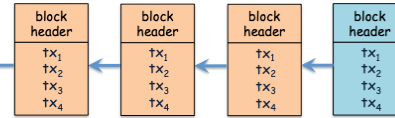
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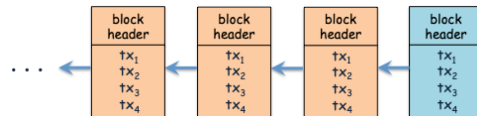


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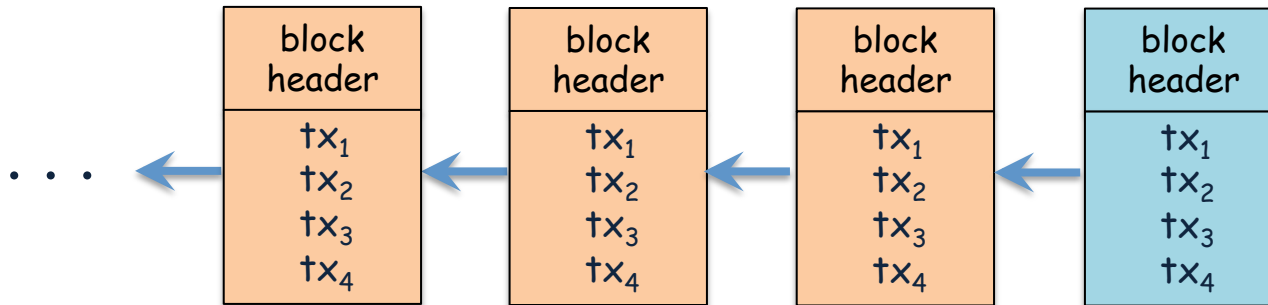


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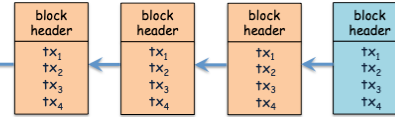
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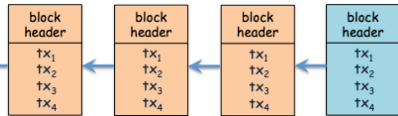
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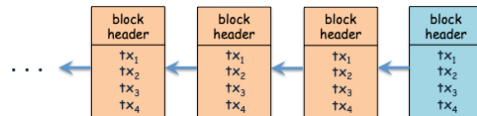
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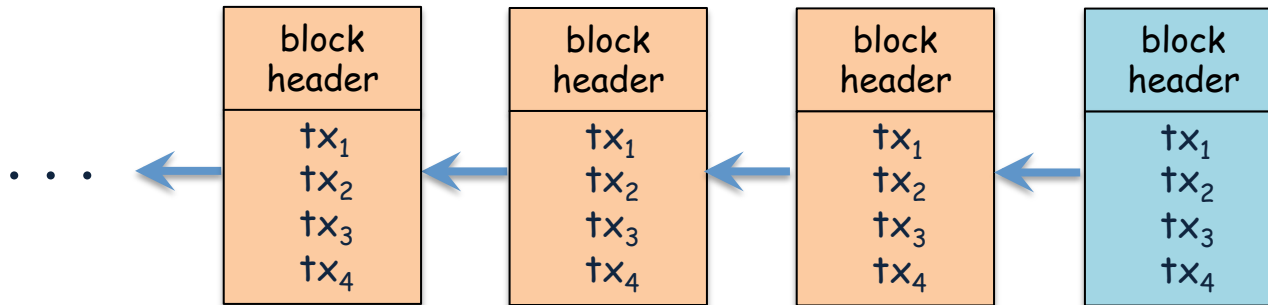
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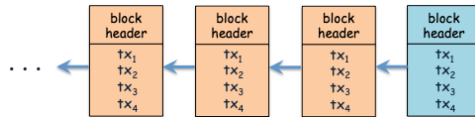
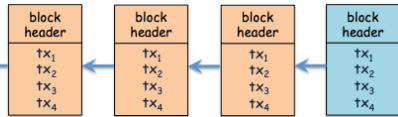
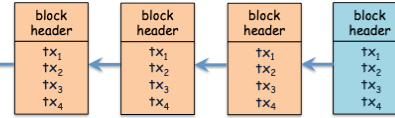
## Proof-of-Work



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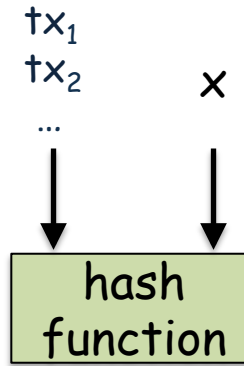
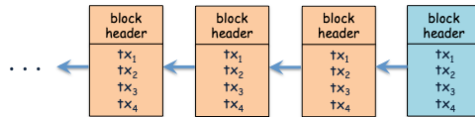
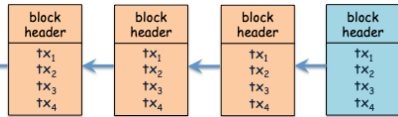
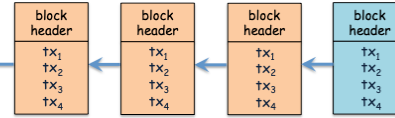


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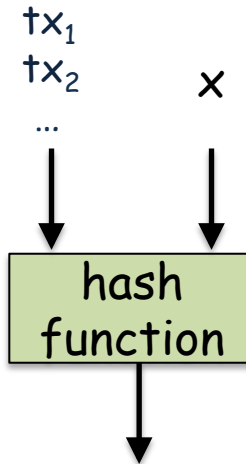
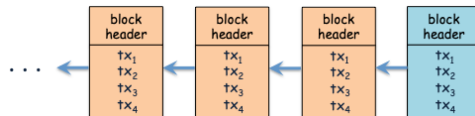
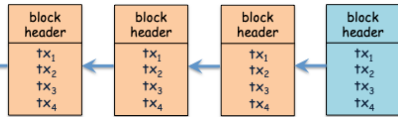
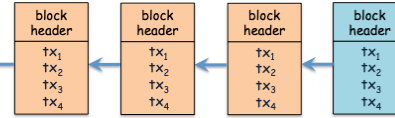


hash  
function

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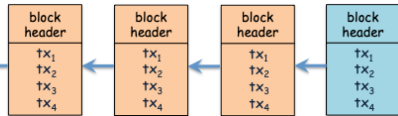
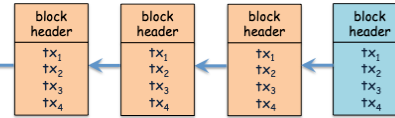


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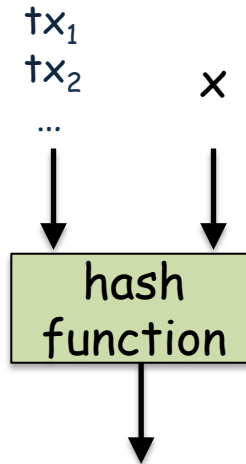
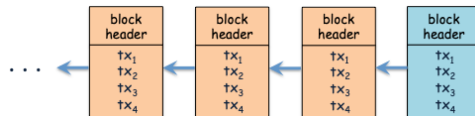


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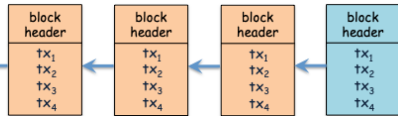
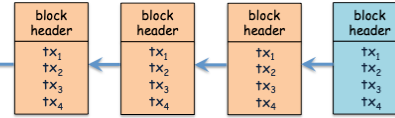


$$H(\text{txs}, x) < T$$

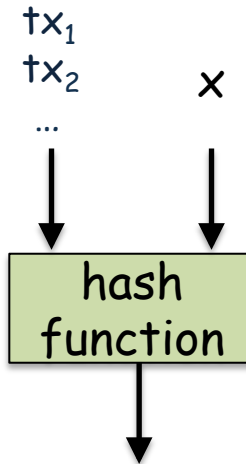
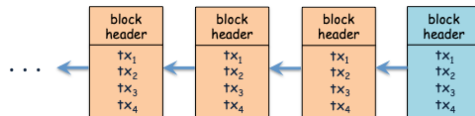


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abc



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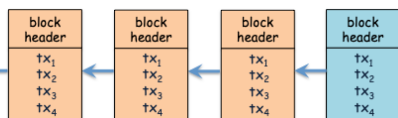
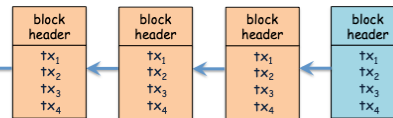
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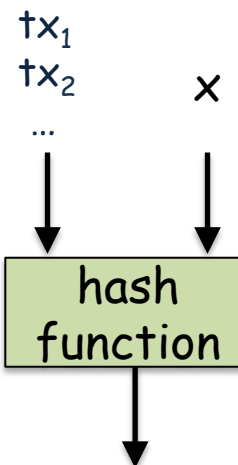
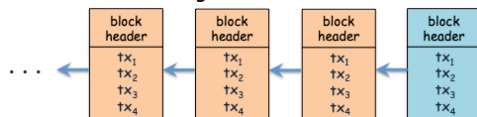
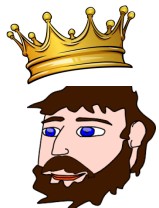
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0f1daffc8a9a3e7866dcd87d34dd16e5

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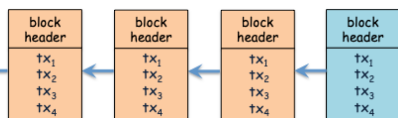
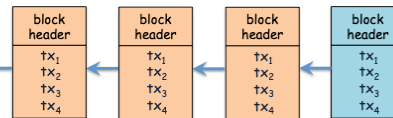
the only way to find x is to try different  
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abC

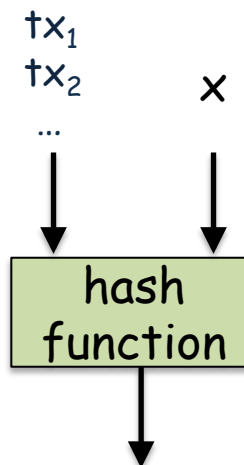
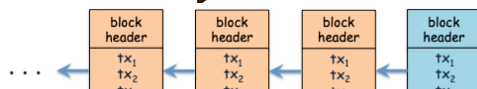


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# Blockchain Consensus Protocols-Nakamoto Consensus Protocol (Proof-of-Work Based)



$$H(\text{txs}, x) < T$$



if i-th player's hash power is  $h_i$ , the probability that i-th player wins the game is  $w_i / \sum w_j$

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abc

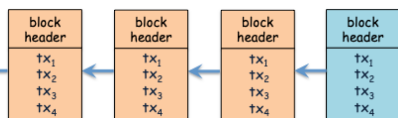
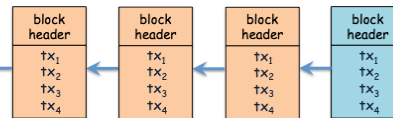
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to keep the block time interval 10 m, the difficulty level  $T$  adjusted in every 2016 blocks (app. 2 weeks)

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abc

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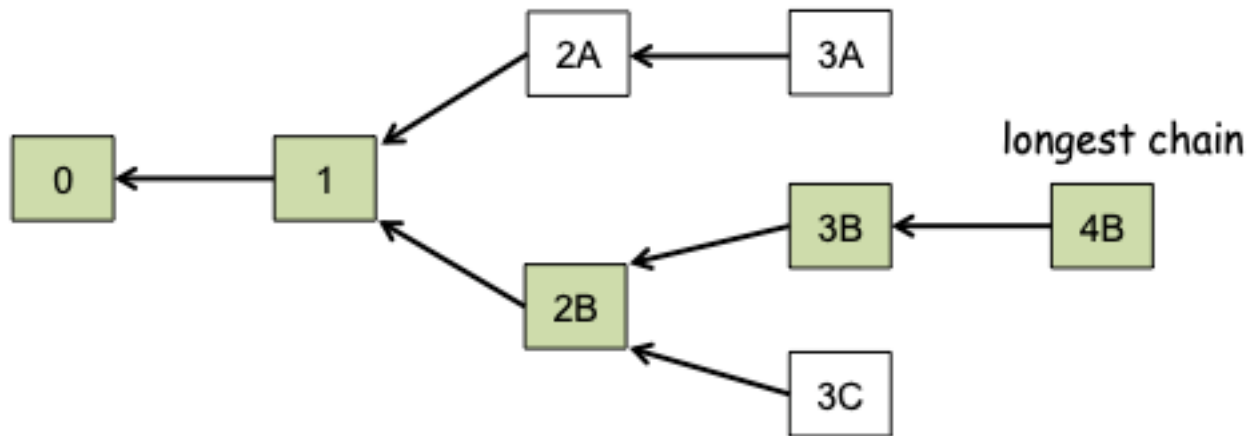
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  - persistence, once a tx recorded more than  $k$  blocks deep in the blockchain of one honest node, then it will be included in every honest node's chain with very high probability
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  - liveness, all txs shared by honest nodes will eventually be placed more than  $k$  blocks deep in the blockchain of an honest node's chain
- they [7] that Nakamoto consensus protocol provides persistence and liveness,
  - (i) if adversary controls minority of the total hashing power in the network,
  - (ii) digital signature scheme unforgeable,
  - (iii) network synchronizes much faster relative to PoW solution rate,

# Blockchain Consensus Protocols-Nakamoto Consensus Protocol (Proof-of-Work Based)

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  - to rewrite the some part of the chain,
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| Name                        | Symbol | Market Cap | Algorithm | Hash Rate    | 1h Attack Cost | NiceHash-able |
|-----------------------------|--------|------------|-----------|--------------|----------------|---------------|
| <a href="#">Bitcoin</a>     | BTC    | \$793.12 B | SHA-256   | 183,643 PH/s | \$1,650,017    | 0%            |
| <a href="#">Ethereum</a>    | ETH    | \$370.60 B | Ethash    | 938 TH/s     | \$1,708,465    | 7%            |
| <a href="#">Litecoin</a>    | LTC    | \$9.44 B   | Scrypt    | 340 TH/s     | \$116,597      | 13%           |
| <a href="#">BitcoinCash</a> | BCH    | \$7.02 B   | SHA-256   | 1,559 PH/s   | \$14,009       | 34%           |

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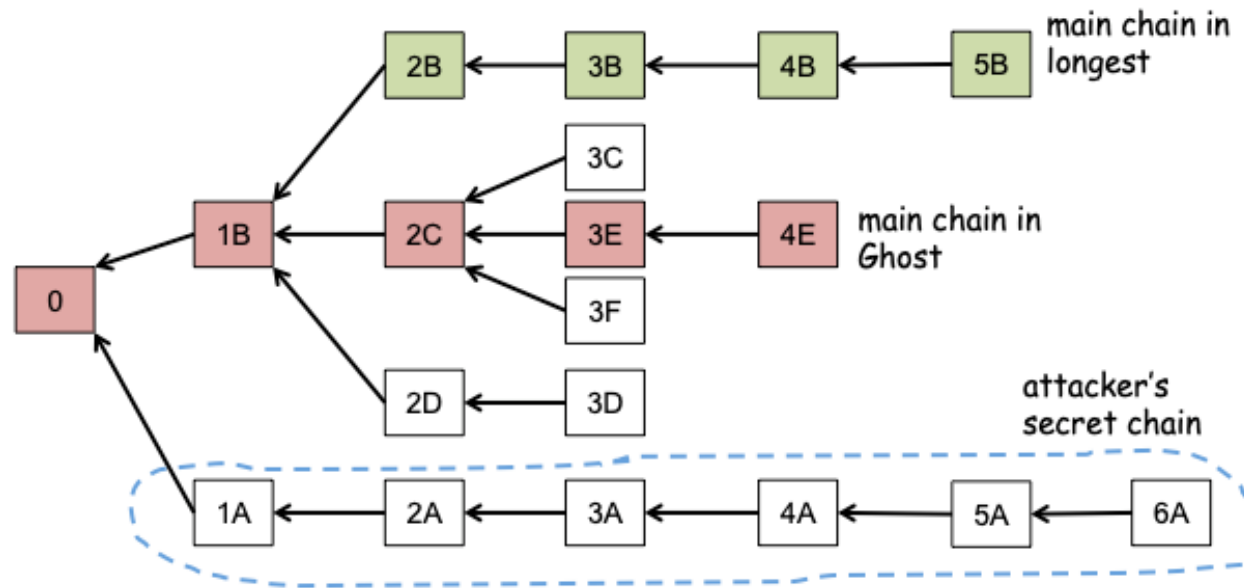
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- transactions per second (tps - throughput) for bitcoin 4-5 (maximum 7) (Visa can process more than 24k [10])
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- Croman et al. [11] showed that when block size increased to 4MB (meaning 26-28 tps) 10% of the nodes would not be able to properly get the newly created blocks (it will reduce the network's effective hash power)
  - if the block size increased to 38MB (meaning 248-250 tps), it will become 50%

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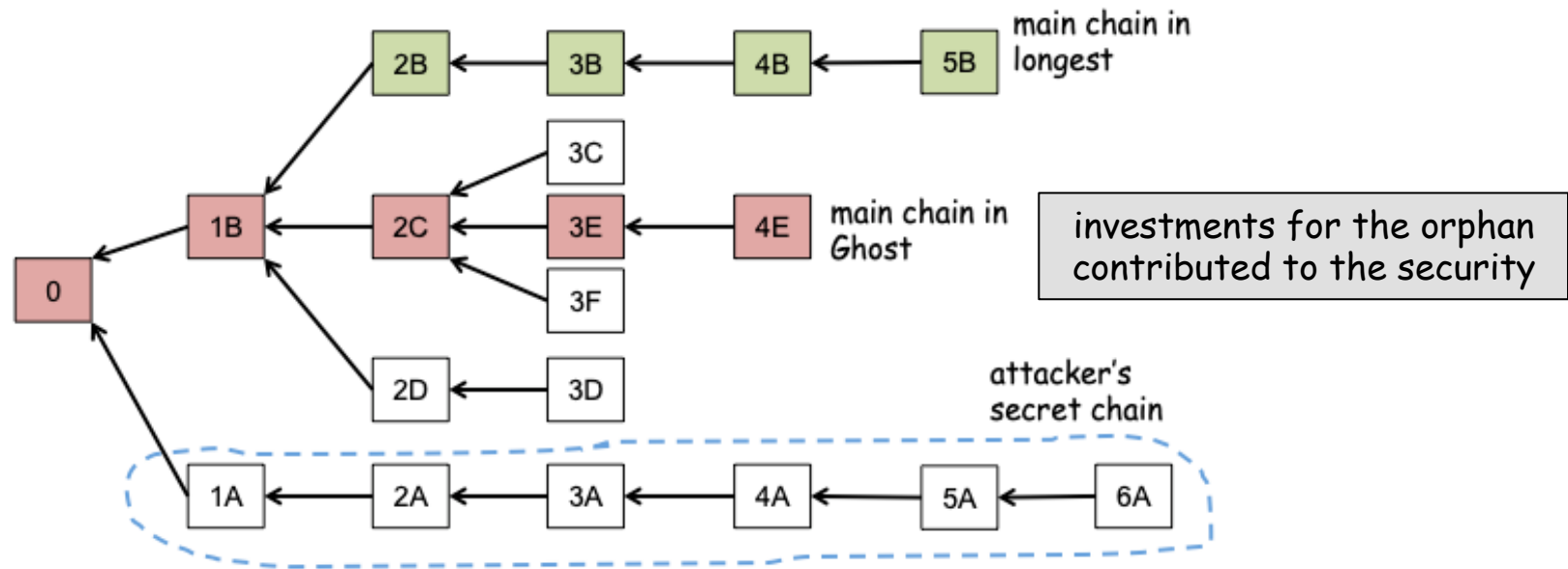
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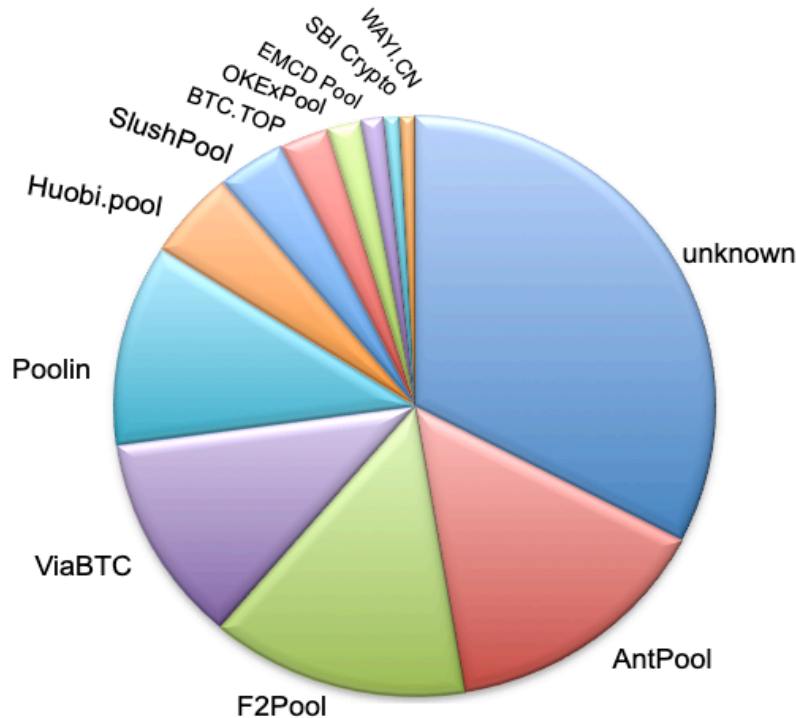
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- they buy special equipments like ASIC ( $\approx$ \$15k-20k) to have an advantage on mining competition
- makes it harder for small players to compete on mining  
(better to join a mining pool)

# Blockchain Consensus Protocols-Nakamoto Consensus Protocol (Proof-of-Work Based)

- a miner can earn \$270k-290k when creating a valid block with bitcoin price \$41k (checked in 4.3.22)
- this reward incentivizes too many people to make investments on mining
- they buy special equipments like ASIC ( $\approx$ \$15k-20k) to have an advantage on mining competition



- makes it harder for small players to compete on mining (better to join a mining pool)

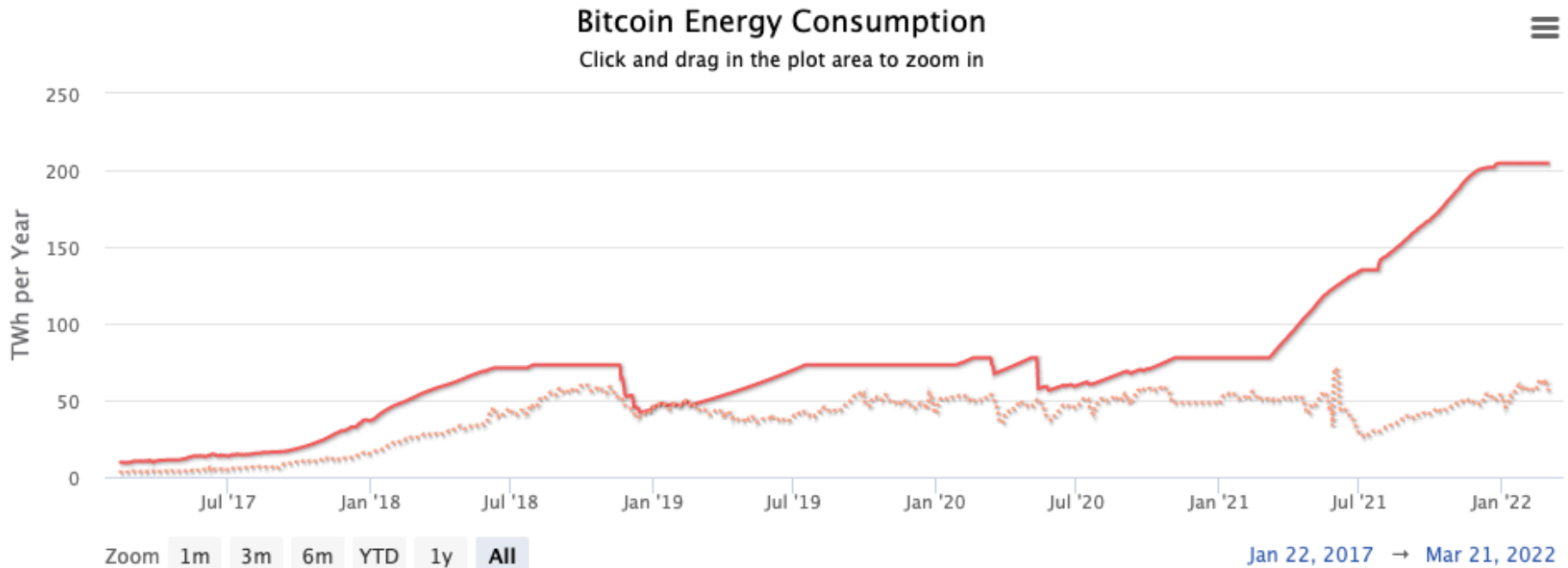
taken from  
blockchain.com  
in 13.06.21

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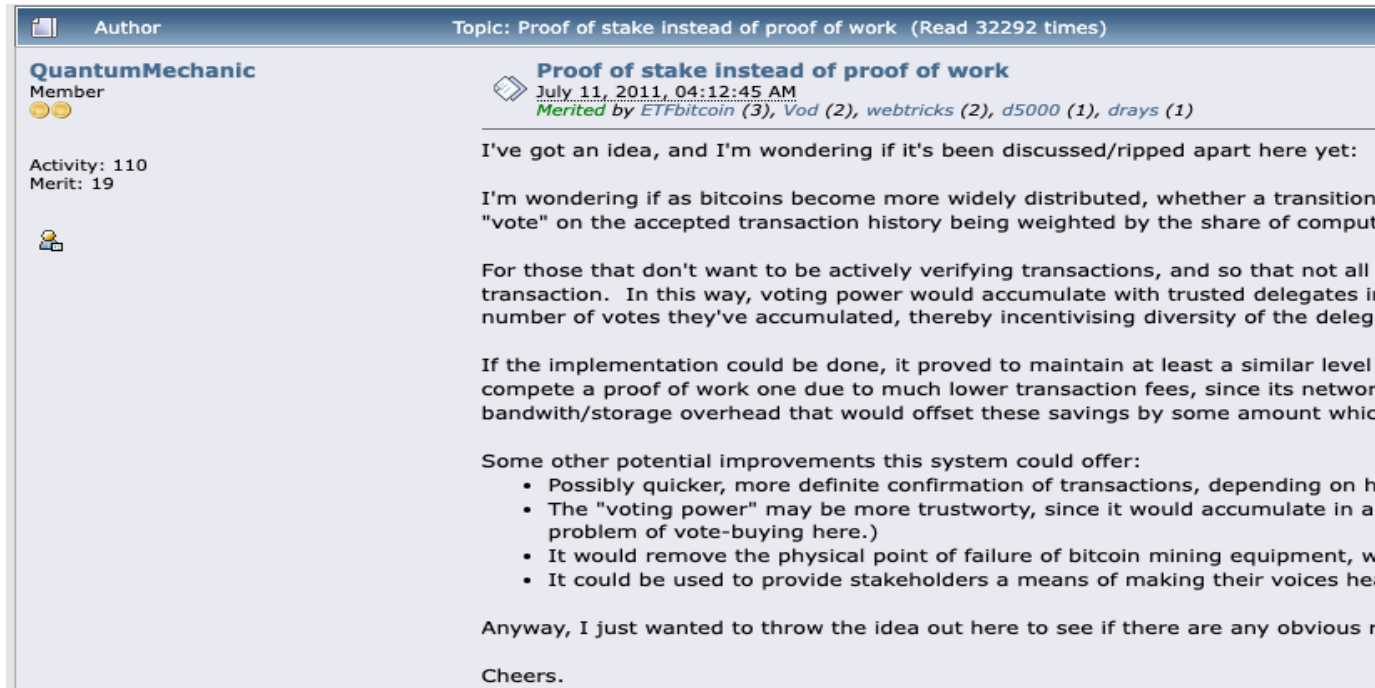


taken from [digiconomist.net](https://digiconomist.net) in 13.06.21

- more than Argentina and Holland

# Blockchain Consensus Protocols - Proof-of-Stake-Based

- the idea first introduced by QuantumMechanic in 2011:  
"instead of your "vote" on the accepted transaction history being weighted by the share of computing resources you bring to the network, it's weighted by the number of bitcoins you can prove you own, using your private keys"



The screenshot shows a forum post by user QuantumMechanic. The post title is "Proof of stake instead of proof of work" and it has been read 32,292 times. The post was made on July 11, 2011, at 04:12:45 AM and has been merited by several users: ETFbitcoin (3), Vod (2), webtricks (2), d5000 (1), and drays (1). The post content discusses the idea of a proof of stake consensus mechanism as an alternative to proof of work, highlighting its potential benefits like lower transaction fees and reduced bandwidth/storage overhead. It also lists some potential improvements such as faster confirmation times and increased trustworthiness.

Author: QuantumMechanic (Member, Activity: 110, Merit: 19)

Topic: Proof of stake instead of proof of work (Read 32292 times)

**Proof of stake instead of proof of work**  
July 11, 2011, 04:12:45 AM  
Merited by ETFbitcoin (3), Vod (2), webtricks (2), d5000 (1), drays (1)

I've got an idea, and I'm wondering if it's been discussed/ripped apart here yet:

I'm wondering if as bitcoins become more widely distributed, whether a transition "vote" on the accepted transaction history being weighted by the share of computing resources you bring to the network, it's weighted by the number of bitcoins you can prove you own, using your private keys

For those that don't want to be actively verifying transactions, and so that not all participants need to verify every transaction. In this way, voting power would accumulate with trusted delegates in proportion to the number of votes they've accumulated, thereby incentivising diversity of the delegates.

If the implementation could be done, it proved to maintain at least a similar level of security to a proof of work one due to much lower transaction fees, since its network bandwidth/storage overhead that would offset these savings by some amount which would be determined by the implementation.

Some other potential improvements this system could offer:

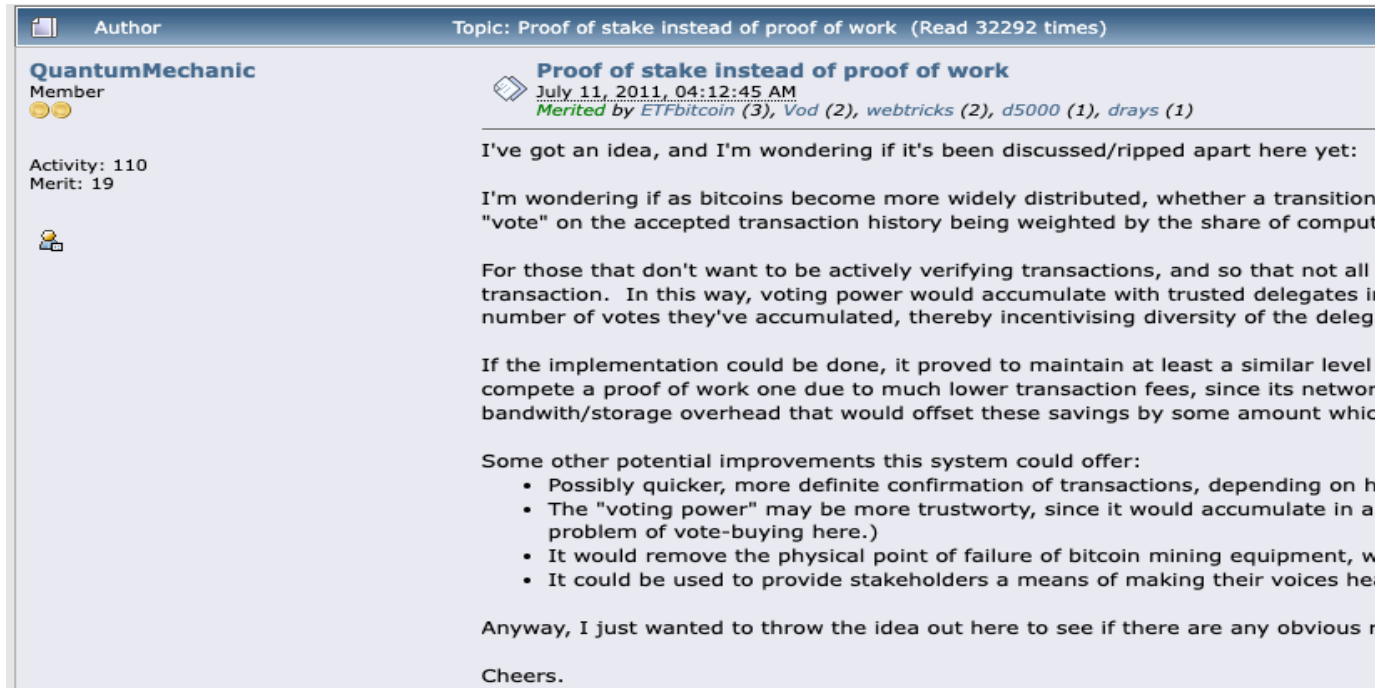
- Possibly quicker, more definite confirmation of transactions, depending on how the system is implemented.
- The "voting power" may be more trustworthy, since it would accumulate in a more distributed manner (no problem of vote-buying here.)
- It would remove the physical point of failure of bitcoin mining equipment, which is a major concern for miners.
- It could be used to provide stakeholders a means of making their voices heard.

Anyway, I just wanted to throw the idea out here to see if there are any obvious reasons why it wouldn't work.

Cheers.

# Blockchain Consensus Protocols - Proof-of-Stake-Based

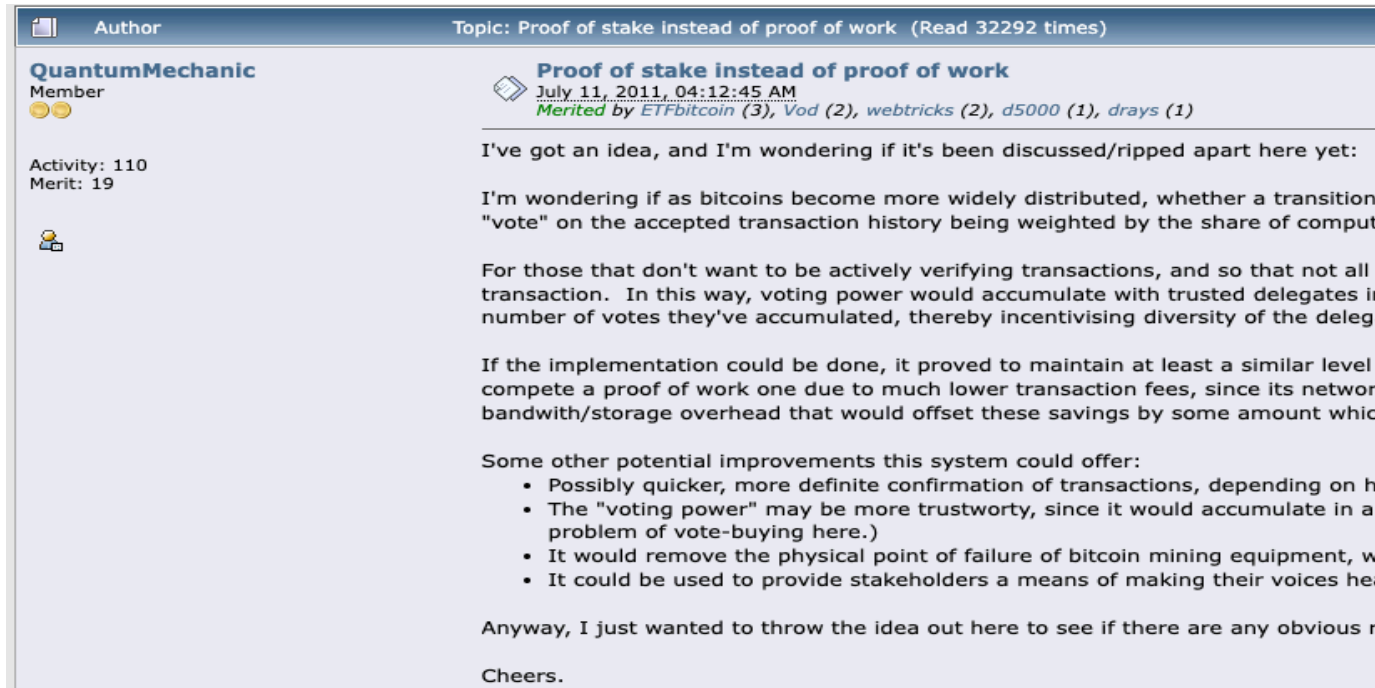
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- the parties who hold the stake in the system are well-suited to maintain the ledger since their stake will diminish in value when the security of the system collapses
- a party who possesses  $p$  fraction of the total amount of coins in circulation will be the leader with the probability  $p$



## Chain of Activity

- introduced by Bentov et al. [13] in 2016
- a pure Proof of Stake protocol that aims to prevent the rational forks by which the only a single stakeholder identity can create the next block

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- introduced by Bentov et al. [13] in 2016
- a pure Proof of Stake protocol that aims to prevent the rational forks by which the only a single stakeholder identity can create the next block
- there are two difficulties associated with pure Proof of Stake system:
  - fair initial distribution of the money supply to the parties
  - network fragility if the nodes are rational

# Blockchain Consensus Protocols - Proof-of-Stake-Based

## Chain of Activity

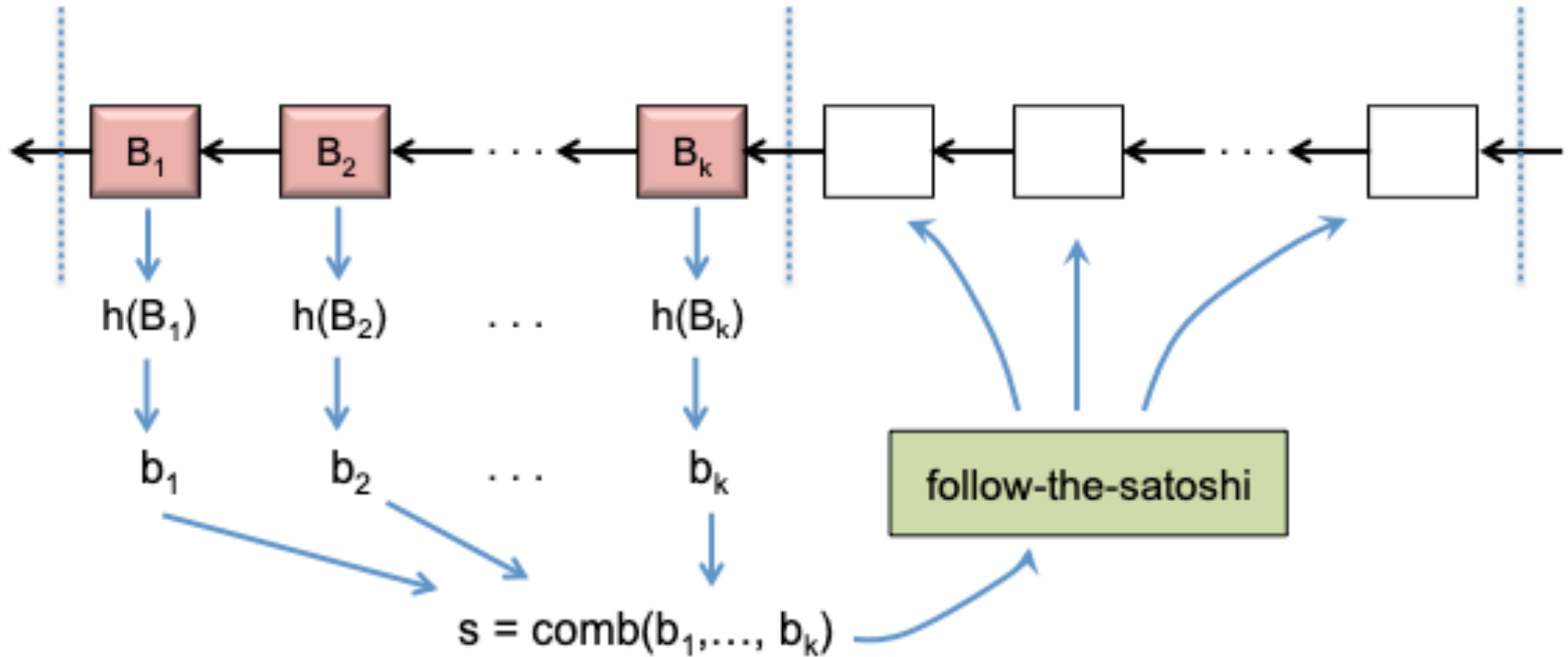
- time is divided into sequence of segments, called epoch
- each epoch is divided into  $L$  discrete unites, called slot
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- the identity of this stakeholder is fixed and publicly known
- the leaders of the current epoch will form a seed as  $SL = \text{comb}(b_1, \dots, b_L)$  where  $b_i = \text{Hash}(B_i)$
- the seed is then used to derive the identities of the next  $L$  stakeholders via 'follow-the-satoshi'

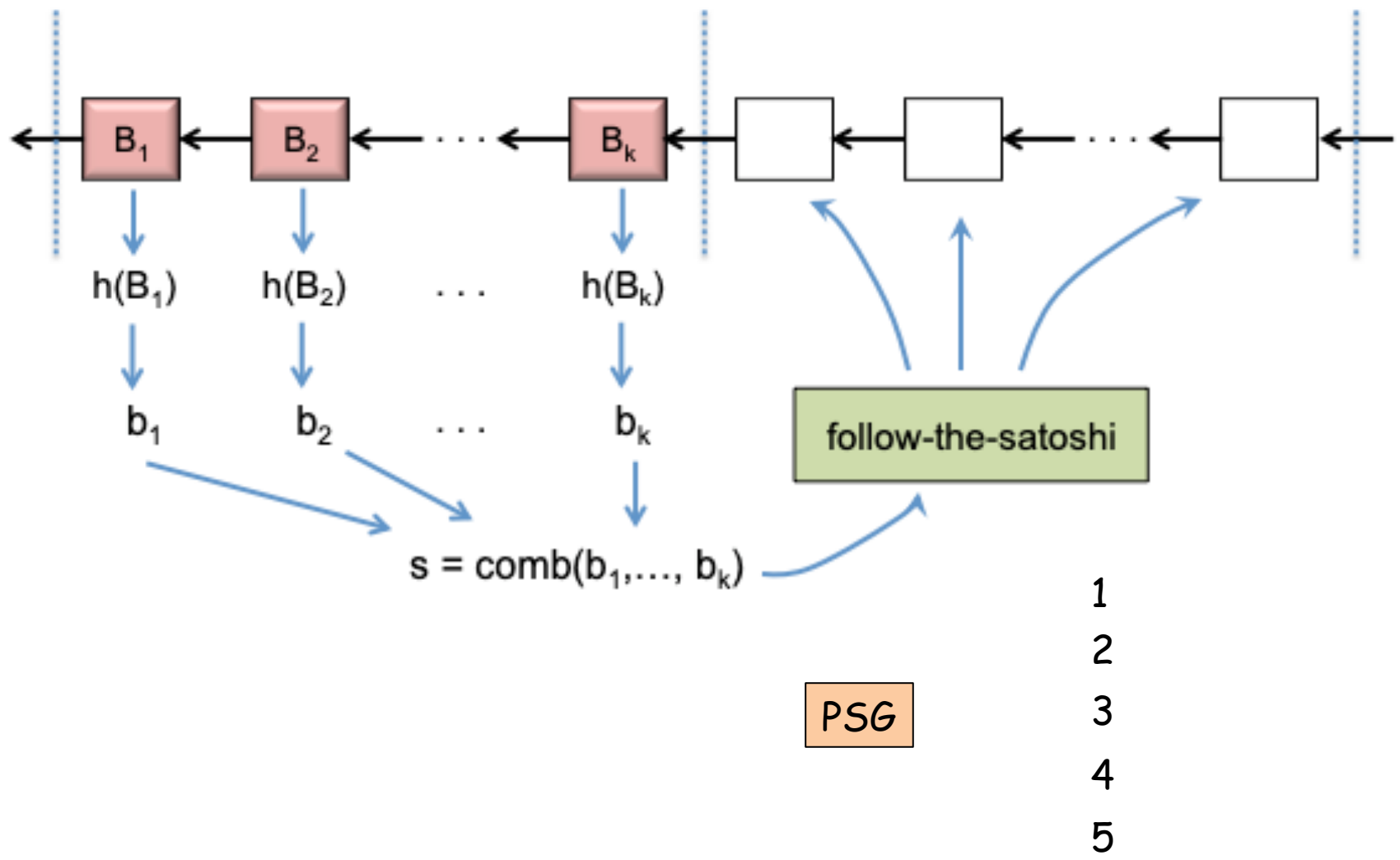
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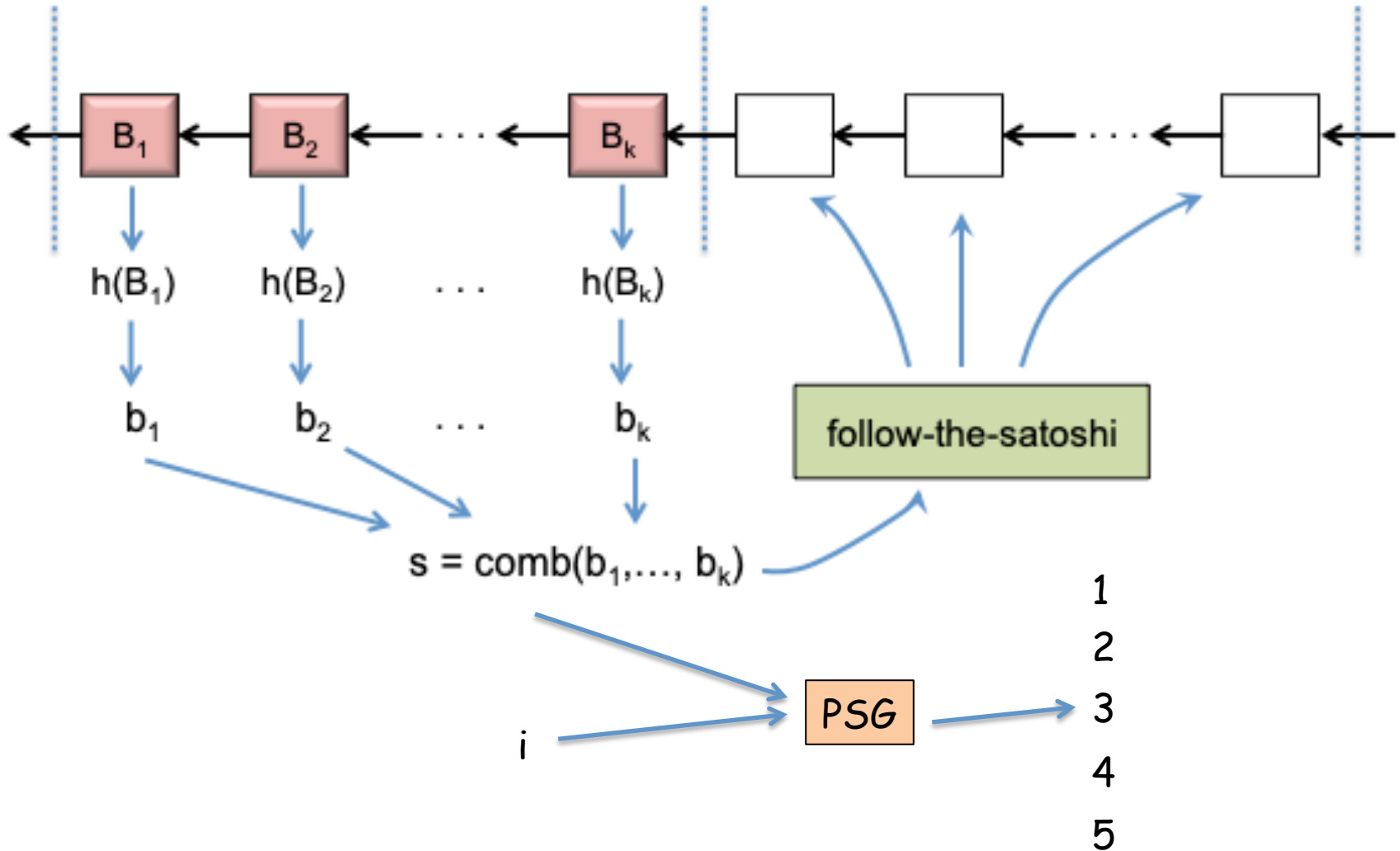
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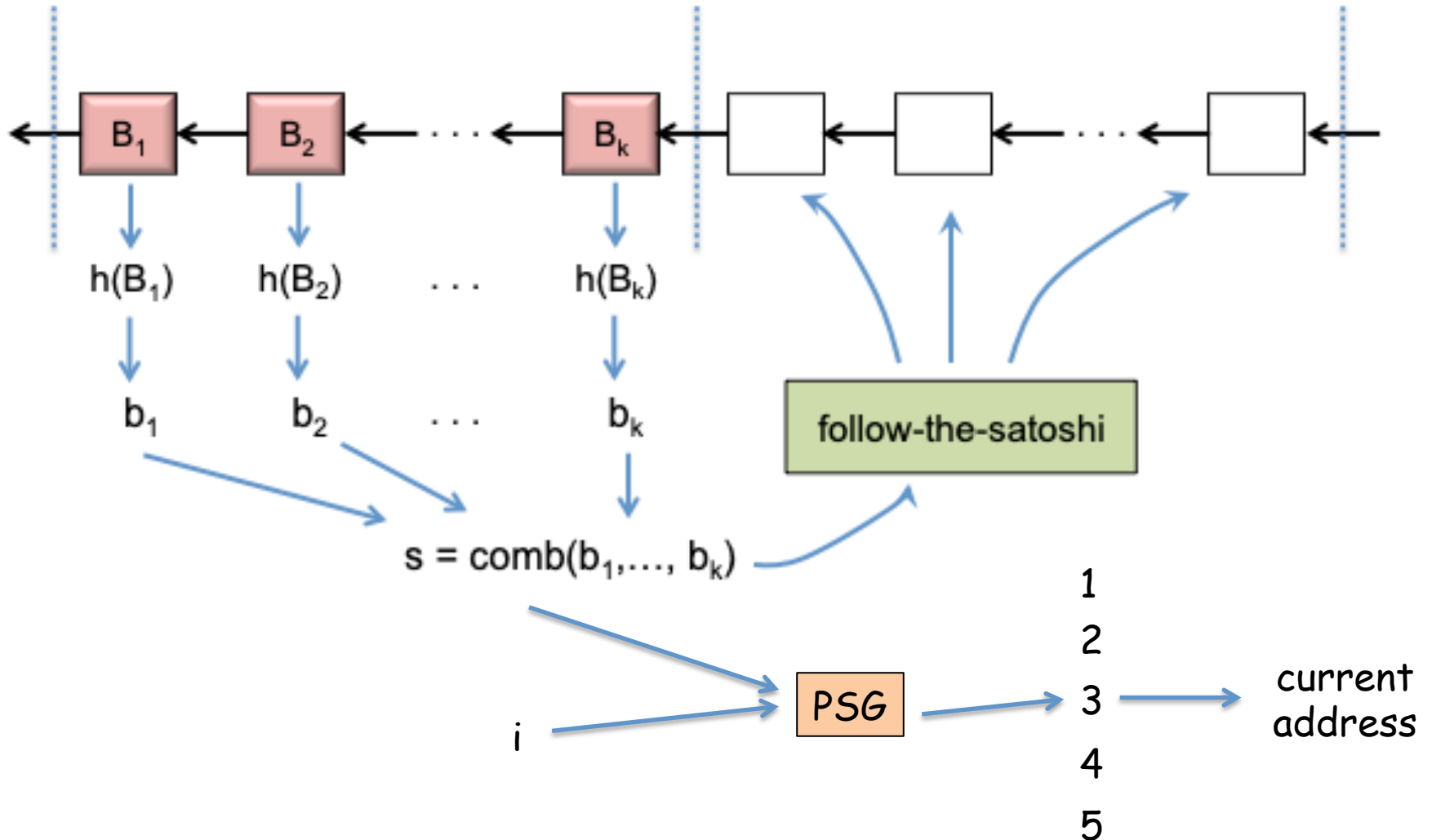
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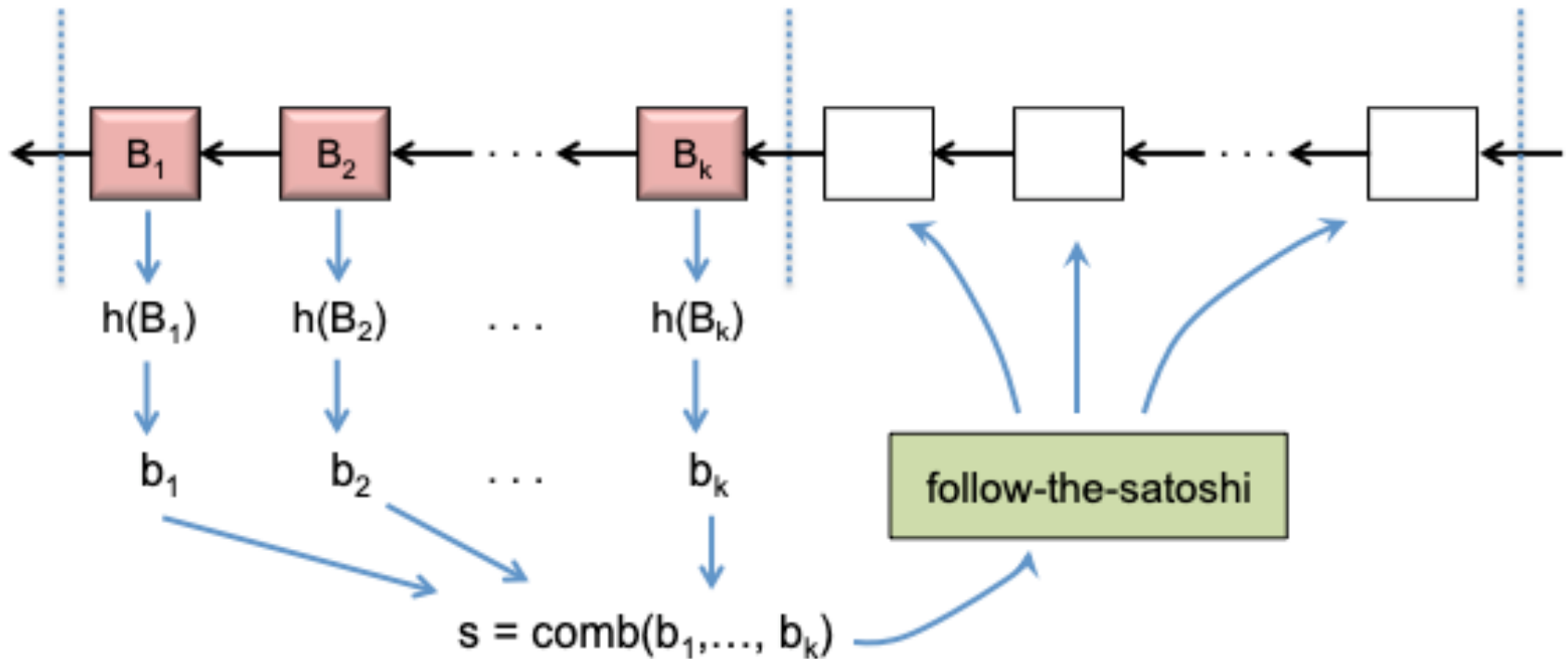
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# Blockchain Consensus Protocols - Proof-of-Stake-Based

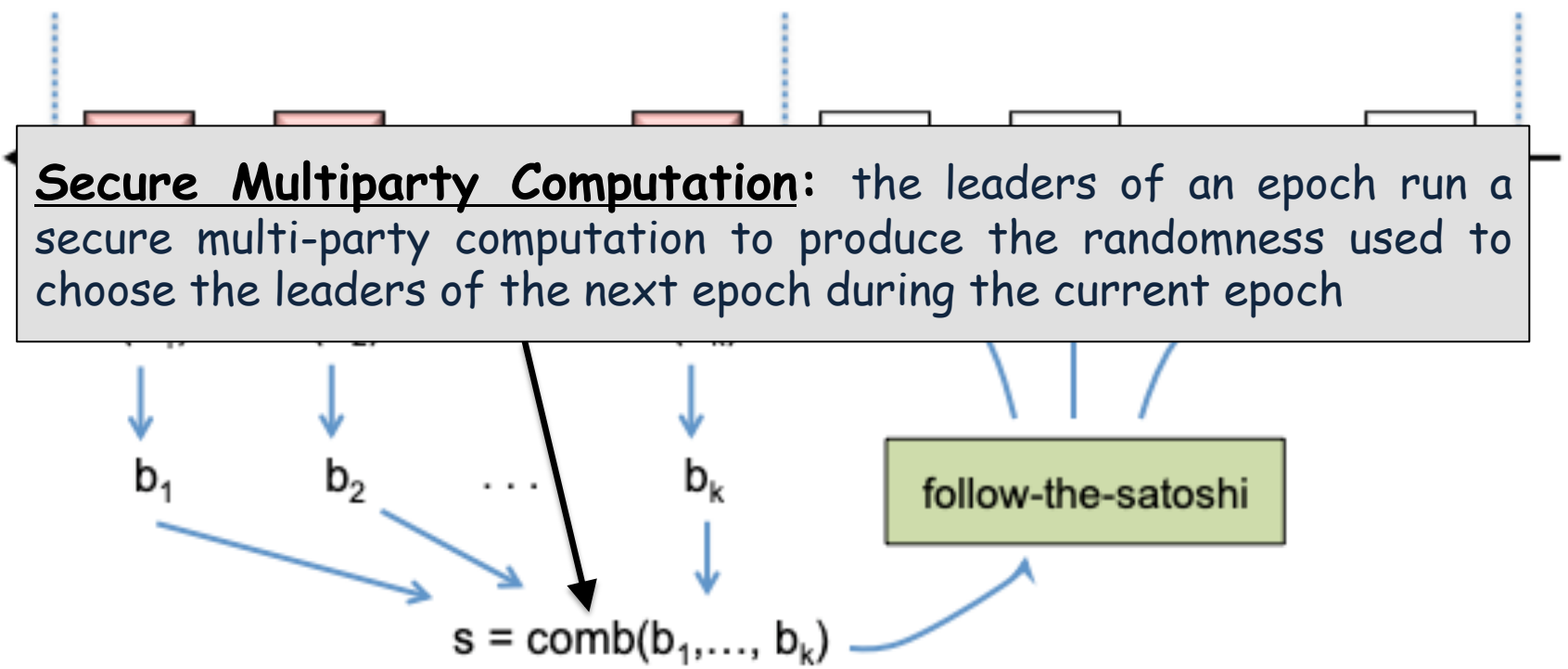
## Ouroboros



- introduced by Aggelos et al. [14] as the first blockchain protocol based on PoS with rigorous security guarantees
- a fundamental problem for PoS is to simulate the leader election process.
- an adversary controlling a set of stakeholders may attempt to simulate the protocol execution trying different sequence of stakeholders participants so that it finds a protocol continuation that favors him

# Blockchain Consensus Protocols - Proof-of-Stake-Based

## Ouroboros



**Secure Multiparty Computation:** the leaders of an epoch run a secure multi-party computation to produce the randomness used to choose the leaders of the next epoch during the current epoch

$b_1$

$b_2$

...

$b_k$

follow-the-satoshi

$s = \text{comb}(b_1, \dots, b_k)$

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# Blockchain Consensus Protocols - Proof-of-Stake-Based

## Delegated PoS

- nodes must be online to issue the blocks when they chosen as slot leaders
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- Delegated PoS enables nodes to delegate their stake to others to represent them in the protocol
- thus, they can contribute their stake to the security of the system without being online

## Delegated PoS

- different than Cardano, in general, at the beginning of each epoch, top K delegates according to the votes they obtain determined and assigned to the time slots in the epoch
- Tron - 27, Lisk - 103, Bitshare - > 1% of total stake
- Cardano - 21600
- PoS-based consensus protocols incentivize nodes to create blocks by giving fees or producing some coin at inflation rate

# Blockchain Consensus Protocols - Proof-of-Stake-Based

- Aggelos et al. [14] showed that Ouroboros consensus protocol provides persistence and liveness,
  - (i) if adversary controls minority of the total stake in the network,
  - (ii) digital signature scheme unforgeable,
  - (iii) network is synchronous
  - (iv) nodes do not remain offline for long periods of time

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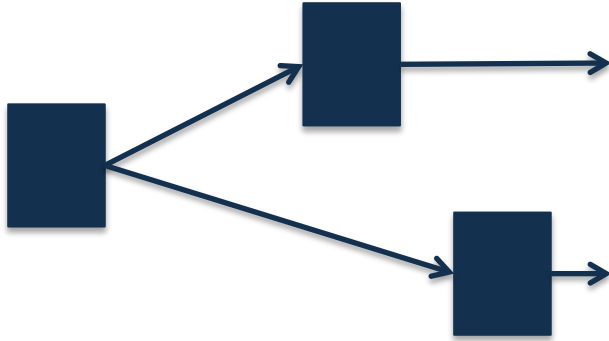
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- lack of formal security proof for most of the protocols (especially for DPoS-based)
- when an attacker gains the control of majority of total stake, it can use it
  - to rewrite the some part of the chain,
  - to damage the network by delaying or censoring some txs
  - to perform double-spending
- similar to Nakamoto statement for PoW, the one acquiring the majority of total stake will use it to improve its gaining not to damage its investments

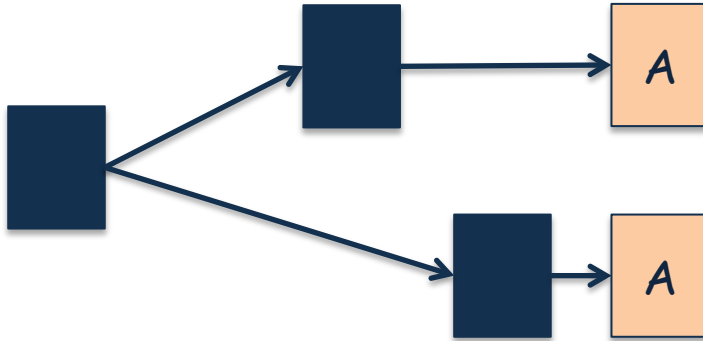


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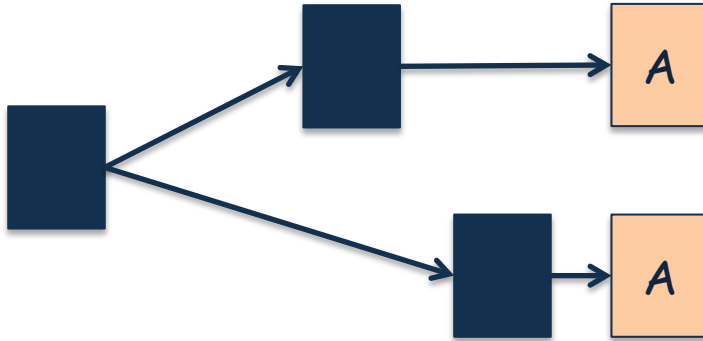
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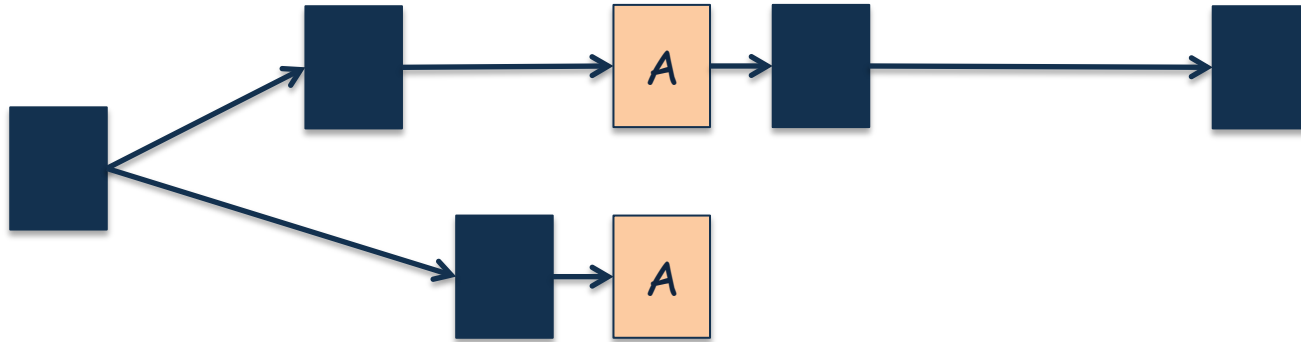
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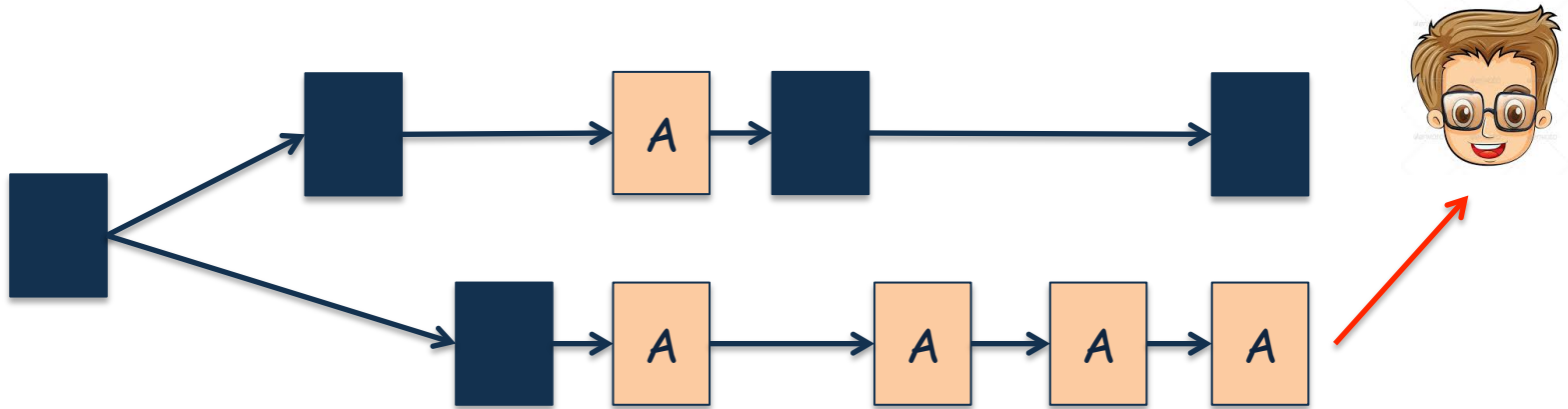
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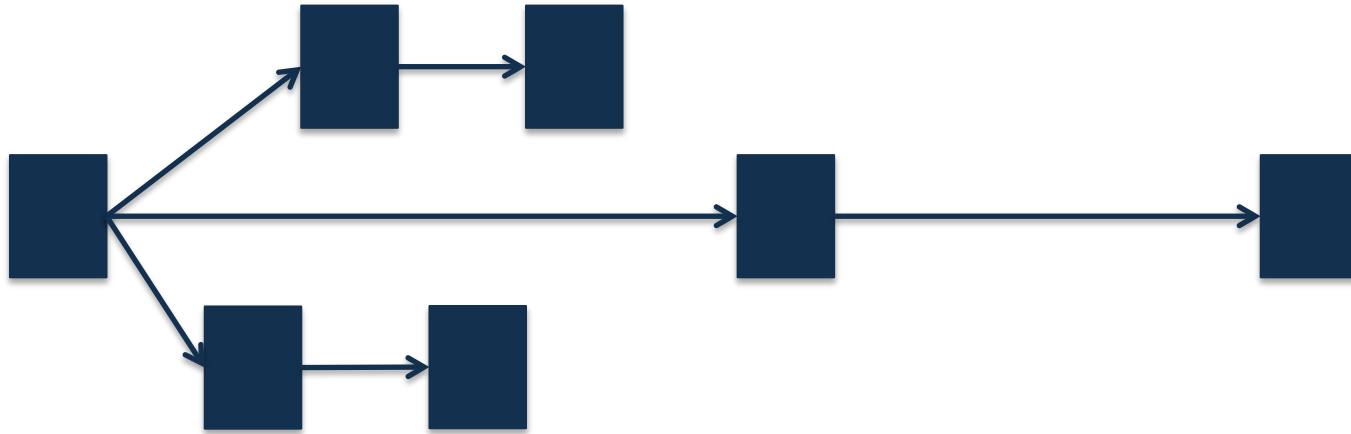
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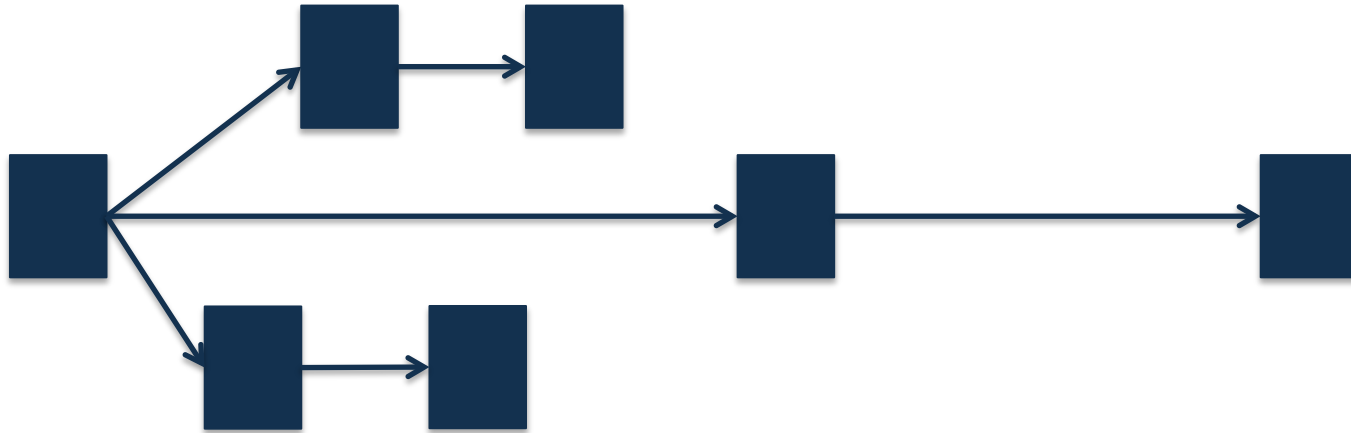
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- Ouroboros proved that this happens almost all the time.

# Blockchain Consensus Protocols - Proof-of-Stake-Based

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# Blockchain Consensus Protocols - Proof-of-Stake-Based

- rich gets richer !
- initial coin distribution ? (73 people for NXT)
- for committee-based PoS, the leaders who issue the blocks determined and shared before each epoch
  - they become targets for some attacks
  - Kerber et al. [15] proposed a protocol that hides the identities of the slot leaders of the next epoch

# Lottery-Based Protocols

| protocol  | leader selection | incentivization      | fault tolerance              | throughput | disadvantage            |
|-----------|------------------|----------------------|------------------------------|------------|-------------------------|
| Bitcoin   | PoW              | fresh coin + fee     | minority of total hash power | 7 tps      | electricity consumption |
| Ghost     | PoW              | fresh coin + fee     | minority of total hash power | 15 tps     | electricity consumption |
| CoA       | PoS              | fee                  | minority of total stake      | ?          | ICD                     |
| Ouroboros | PoS              | fee                  | minority of total stake      | 257 tps    | ICD                     |
| Tron      | DPoS             | fee + inflation rate | minority of total stake      | 2000 tps   | ICD                     |