ANKARA UNIVERSITY COM364 AUTOMATA THEORY

Week 12

Example Questions

Kurtuluş KÜLLÜ

Recall the CFG G_4 that we gave in Example 2.4. For convenience, let's rename its variables with single letters as follows.

 $\begin{array}{l} E \rightarrow E + T \mid T \\ T \rightarrow T \times F \mid F \\ F \rightarrow (E) \mid a \end{array}$

Give parse trees and derivations for each string.

a.	a	c.	a+a+a
b.	a+a	d.	((a))

COM364	M364
---------------	------

Give a context-free grammar (CFG) that generate the following language. The alphabet is {0,1}.

{w| the length of w is odd}

Let $G = (V, \Sigma, R, S)$ be the following grammar. $V = \{S, T, U\}; \Sigma = \{0, \#\};$ and R is the set of rules:

 $S \rightarrow TT \mid U$ $T \rightarrow 0T \mid T0 \mid \#$ $U \rightarrow 0U00 \mid \#$

- **a.** Describe L(G) in English.
- **b.** Prove that L(G) is not regular.

Convert the following CFG into an equivalent CFG in Chomsky normal form, using the procedure given in Theorem 2.9.

 $\begin{array}{c} A \rightarrow BAB \mid B \mid \varepsilon \\ B \rightarrow 00 \mid \varepsilon \end{array}$

Consider the following CFG G:

 $\begin{array}{ccc} S \ensuremath{\:\rightarrow\:} SS \mid T \ T \ensuremath{\:\rightarrow\:} \mathbf{a}T\mathbf{b} \mid \mathbf{a}\mathbf{b} \end{array}$

Describe L(G) and show that G is ambiguous. Give an unambiguous grammar H where L(H) = L(G) and sketch a proof that H is unambiguous.

Design pushdown automata (PDAs) that recognize the languages given in questions 5 and 3.